

Scholars' Mine

Masters Theses

Student Theses and Dissertations

1984

CIEGEN: A system for testing knowledge base compilation heuristics on a microcomputer

Jayne D. Ward

Follow this and additional works at: https://scholarsmine.mst.edu/masters_theses

Part of the Computer Sciences Commons Department:

Recommended Citation

Ward, Jayne D., "CIEGEN: A system for testing knowledge base compilation heuristics on a microcomputer" (1984). *Masters Theses*. 4544. https://scholarsmine.mst.edu/masters_theses/4544

This thesis is brought to you by Scholars' Mine, a service of the Missouri S&T Library and Learning Resources. This work is protected by U. S. Copyright Law. Unauthorized use including reproduction for redistribution requires the permission of the copyright holder. For more information, please contact scholarsmine@mst.edu.



CIEGEN: A SYSTEM FOR TESTING KNOWLEDGE BASE COMPILATION HEURISTICS ON A MICROCOMPUTER

BY

JAYNE D. WARD, 1960-

A THESIS

Presented to the Faculty of the Graduate School of the

UNIVERSITY OF MISSOURI-ROLLA

In Partial Fulfillment of the Requirements for the Degree

MASTER OF SCIENCE IN COMPUTER SCIENCE

1984

Approved by

Billy Lillett (Advisor) John Klank

ABSTRACT

The expert system has proven itself to be a valuable aid in diagnosing and treating problems in domains requiring expertise. The commercial world has been alerted to this fact and the thrust is to make the expert system portable and available on small computers.

The goal of this research has been to lay the groundwork for a domain independant expert system builder on a microcomputer. The result of this effort was CIEGEN, a system consisting of a rule compiler, inference engine, and rule generator developed on the IBM PC. It is domain independant, responsible for transforming a knowledge base of rules into heuristic based decision trees, and capable of performing backward chaining consultations.

The system is also heuristic independant, allowing a knowledge base to be compiled by different heuristics and compared using the log created by the inference engine. A subgoal of the development of CIEGEN has been to study the heuristics used to compile a knowledge base because the efficiency of the expert system is based on the intelligence of the heuristic. The heuristic used by EMYCIN was implemented and compared with a heuristic developed by the author. For the six types of knowledge based generated by CIEGEN's rule generator, EMYCIN's heuristic, on the average, executed more quickly.

ii

ACKNOWLEDGEMENT

The author would like to recognize and express her appreciation to all of the people who contributed to the development of this thesis. Dr. Bill E. Gillett and Dr. Arlan R. DeKock have maintained enthusiasm and been very supportive throughout all phases of this research. They both have contributed a considerable amount of time and effort discussing and developing ideas for this thesis. Dr. Ray B. Kluczny reviewed this work in a timely manner; his comments and criticisms were helpful. Alan D. Christiansen was very helpful in the intial stages of CIEGEN. Finally, a special thanks is due to my family, whose fervent and continuous support enabled the completion of this research.

iii

TABLE OF CONTENTS

				Page
ABSTRACT				.ii
LIST OF 2	TABLES	5		• • V
I.	INTRO	DUCT	[ON	1
II.	RELAT	red Li	TERATURE	4
	А.	RULE	COMPILATION	4
	в.	EXPER	RT SYSTEM BUILDING TOOLS	6
III.	SYSTI	EM DES	SIGN	7
	Α.	RULE	BUILDER	7
		1.	RULES	8
		2.	CYCLING	9
			a. FIRST ORDER CYCLING	9
			b. SECOND ORDER CYCLING	.12
		3.	RANDOM NUMBER GENERATOR	.12
		4.	PARAMETERS	.14
	в.	RULE	COMPILER	.15
		1.	ALGORITHM	.16
		2.	DATA STRUCTURES	.17
			a. DECISION TREES	.17
			b. DECISION TREE LISTS	.18
	с.	HEURI	STICS FOR ANTECEDENT SELECTION	.18
		1.	MOST OFTEN OCCURRING	.19
		2.	MINIMUM AVERAGE ANTECEDENT	.19
	D.	INFE	RENCE ENGINE	.21
IV.	EXAMI	PLE CO	DNSULTATION	.22

۷.	SUMMARY	OF RESULTS AND CON	ICLUSIONS24
VI.	FURTHER	RESEARCH SUGGESTIC	DNS32
BIBLIOGR	АРНҮ	•••••••••••••••••	
VITA	••••		
APPENDIX	A	• • • • • • • • • • • • • • • • • • • •	

LIST OF TABLES

Table		page
I.	CONSULTATION	RESULTS

I. INTRODUCTION

The field of Artificial Intelligence (AI) is a subfield of Computer Science and was founded for the purpose of creating systems that could acquire and apply knowledge to problem solving. Researchers discovered, upon attempting to write a program that incorporated intelligence, that very little was understood about the acquisition and storage of knowledge in the human brain. Therefore, one goal of the field is to uncover underlying mechanisms of human intelligence through the developments in AI. The diversity of what constitues intelligence has led to the creation of areas of research in, for example, natural language understanding [1, 2, 3], robotics [4], vision [5], and knowledge engineering [6].

Knowledge engineering, specifically the building of expert systems, is the area of AI that directly relates to this research. The difference between an expert system and a traditional computer program is demonstrated by the types of problems it solves. An expert system typically solves ill structured problems, or problems with incomplete data. It does not solve problems requiring "number crunching" or problems that can be solved by plugging values into a formula as some of the traditional data processing programs.

Expert systems solve problems that require expertise in the areas of, for example, diagnosis, interpretation of data, monitoring, repair, and design. Expertise implies a combination of heuristics or rules of thumb, textbook knowledge, and reasoning. Expert systems use self-knowledge or meta-knowledge enabling them to reason about their solutions and they typically have explanation facilities to justify their solutions. The expert system asks the user for information it needs to solve the problem and allows the user to ask it questions, just as the human expert operates.

Expert systems have been built in many diverse domains. Three examples of expert systems and their domains are: DENDRAL determines molecular structures of unknown compounds [7, 8, 9], MYCIN diagnoses infectious diseases [9, 10, 11], and PROSPECTOR gives advice on finding ore deposits from geological data [12].

The most desirable way to build an expert system is through a domain independant tool insuring the separation of the knowledge and the control structures. A desirable feature of this tool is the ability to build an expert system on a microcomputer making it convenient for most people in the commercial field.

This paper describes a system, CIEGEN, which was designed with an emphasis on the features of domain independance, portability, and efficiency. CIEGEN, consisting of a rule compiler, inference engine, and rule generator, aids the user in building an efficient knowledge base on the IBM PC.

The knowledge base contains the expertise used to solve problems in a particular domain such as medicine or geology. A popular representation for the expertise is rules of the

form: if (condition) then (action), because of their modularity, representing single "chunks" of knowledge. Other representations of knowledge include frames and semantic nets, but these will not be addressed in this paper.

In CIEGEN, these rules are transformed into a representation that will execute more efficiently by a process called rule compilation. Rules are compiled into decision trees which effectively allow parallel execution of several rules at once. The antecedent (condition) chosen as a branch of the decision tree is selected by a heuristic. Since the heuristic is responsible for the efficiency of the knowledge base it is important to be able to compare the results of different heuristics. This need is provided by CIEGEN.

Literature related to the compilation process and expert systems are reviewed in Chapter II. Chapter III describes CIEGEN's rule generator, compiler, and inference engine. Demonstrating the usefulness of CIEGEN, two heuristics for rule compilation are compared in this paper. One heuristics is currently in use by an expert system builder, EMYCIN, and the other heuristic was developed by the author. Both of these are described in Chapter III. An example of the compilation and consultations are given in Chapter IV. The results of this research are summarized in Chapter V and suggestions for further research are given in Chapter VI.

II. RELATED LITERATURE

A. RULE COMPILATION

Many researchers agree, that as larger knowledge bases are required for expert systems, techniques such as rule compilation will become mandatory [9]. Heuristic compilation is equivalent to establishing the search strategy. The exception to the equivalency is the metaknowledge that may be applied at execution time to alter the search path. As knowledge bases become large it is necessary to perform intelligent searches to keep costs from exponentially increasing. It is also important to eliminate redundancy in testing of similar patterns in rules which constitutes the fundamental compilation algorithm.

Researchers working with the EMYCIN system conducted a study comparing consultation times of expert systems using compiled knowledge bases with consultation times of intepreted knowledge bases. Results showed that the interquestion or "think time" was cut close to half for the systems PUFF, SACON, and MYCIN [13].

These systems were backward chaining or goal directed systems. This means that a goal is established, rules concluding about this goal are gathered, and the conditions in these rules become the new subgoals. This process continues until all conditions in a rule are known, the needed actions are executed. To compile a backward chaining knowledge base means that all rules concluding about a single parameter are located in one decision tree. Rather than searching the knowledge base for all rules concluding about a particular parameter, the inference engine simply travels down the branches of the decision tree.

Compilation has also proven effective in the data driven or forward chaining systems. A data driven system begins with known values in what is called working memory, matches the left hand side of the rules with the known values and draws conclusions from those rules. The conclusions enter working memory and cause other rules to be candidates for execution. The problem with these systems is that with large knowledge bases, the matching process is very slow because it has to repeatedly check elements in working memory. Data driven systems have been reported to spend over nine-tenths of their run time performing the matches [9].

The most common attempt at improving the efficiency of these systems, known as production systems, has been by combining indexing with interpretations of the left hand sides. A successful implementation of the compilation process has been developed by Forgy [14] which is the Rete Match Algorithm. The compiler exploits the properties of similar conditions and the fact that individual productions only change a few facts in memory. Forgy showed, through his studies, that by compiling the productions, the execution time was cut by several orders of magnitude.

B. EXPERT SYSTEM BUILDING TOOLS

An expert system building tool is a domain independant system allowing the development of expert systems in several domains. They prompt the knowledge engineer or expert for rules (knowledge), parameters (goals), and definitions of the parameters. They provide the control structure for the expert system, which includes the inference engine and compiler (if this technique is used).

The level of interaction between the user and the system varies among different systems. For example, EMYCIN prompts its user with a terse Abbreviated Rule Language which is a cross between LISP and English. Teiresias [15, 16] interacts with its user in reasonable English. One limitation of this system, due to the difficulty of parsing English, is the assumption that a dictionary and knowledge base have already been established and the user is merely editing the knowledge base. Other systems such as KAS [9] are not as versatile and were developed specifically for use with an expert system (PROSPECTOR).

III. SYSTEM DESIGN

CIEGEN is a system consisting of a compiler, inference engine, and automatic rule generator. The knowledge base used for this research was generic in the sense that the rules consist of arbitrary alphabetic letters with no particular meaning assigned to them. The reason for using a generic knowledge base as opposed to a particular domain was to permit clearer recognition of the results from the compiler.

CIEGEN was developed on an IBM PC and written in IQLISP [17]. Since the IQLISP environment occupies approximately 106k bytes of RAM, the machine should be equipped with at least 256k of RAM.

In order to conduct this research a total of six packages were needed. They include the Rule Builder, the Rule Compiler, Inference Engine, one of the heuristics for compilation: Heurisl (Most_Often_Occurring), Heuris2 (Minimum Average Antecedent), Rules which is a general utility package used by all other packages, and FLOAT which is a package to enable real arithmetic.

A. RULE BUILDER

The decision to mechanically generate rules was made to allow control over certain parameters describing the knowledge base. Some of those parameters are: the number of rules, the number of unique consequents, the number of

antecedents per rule and the number of knowns per rule. Another factor leading to the decision to automatically generate rules was to be able to maintain randomness and to avoid creating rules that would favor one heuristic over another.

1. Rules

The rule (A B C ==> D) is read as:

[If A and If B and If C are true then conclude D is true], so that if A or B or C is false then no conclusion is made. Similarly, a rule involving a "not" (where # = "not") such as (#A B C ==> D) reads "not" A and B and C in order to conclude D is true.

In CIEGEN, the structure of a rule is a LISP list ((A B C) D) where the CAR of the rule is the list of antecedents and the CADR of the rule is the conclusion.

Each rule will have only one consequent and its certainty factor will be one (assuming no probability is involved). This first constraint could be changed by allowing the conclusion to be a list rather than an atom and the latter by making the certainty factor a property of the rule. However, for the purpose of studying different heuristics for compiling rules, the former structure proved sufficient.

As rules are generated, certain information about them is stored. For example, the total number of antecedents and the number of knowns are stored in an array called ANT_INFO

to be used for examination and for the Minimum Average Antecedent heuristic. Also stored is each upper case letter used in a rule. An upper case letter represents something that is unbound, it will have a set of rules concluding about it. For example, suppose the first two rules in the knowledge base are [((P G L i) A) ((F e d M) A)]. The first and second rows of ANT_INFO would be as follows:

Al4 PGL

A 2 4 F M , where the ASCII values of the the letters are used because the array is all integer (A = 65, P = 80, G = 71, L = 76, F = 70, M = 77). The lower case letters are used to represent known values. They are initially given to be true or false corresponding to information asked of a user in a typical expert system consultation. Therefore, they do not have rules concluding about them.

2. Cycling

a. <u>First Order Cycling</u> is something that is expected to happen if nothing is done to prevent it. First order cycling is demonstrated as follows:

B C D ==> AA E F ==> B .

So in order to conclude A the value of B is needed, but in order to conclude B the value of A is required. This problem was initially eliminated by allowing only letters that appear later in the alphabet be candidates for antecedents.

For example, to conclude A, B - Z are candidates, and to conclude B, C - Z are candidates.

Lower case letters a - e are also candidates to be chosen as antecedents. One reason for the lower case letters is to increase the size of the antecedent bucket, providing something to choose from when antecedents are needed for conclusions later in the alphabet. For example, rules to conclude about X have only Y and Z, Y has only Z, and Z has no letters to choose from if the lower case letters are not included.

After analyzing the resulting rules generated by this method, it was felt that part of the randomness was lost by restricting the bucket of candidates. Letters later in the alphabet appeared often in the earlier rules which favored one heuristic over another.

Thus, a decision was made to allow cycling to occur by keeping the antecedent bucket the same throughout the rule generation. After all rules had been generated, a check for first order cycling was made and corresponding rules were eliminated from the knowledge base. To demonstrate this, suppose we have the rule (Z C P G ==> H). This method says to check the rules concluding about C to see if H was used as an antecedent. If it was, as in the rule (X \underline{H} S D ==> C), then a cycle exists and the rule (Z C P G ==> H) is removed from the knowledge base. If H did not appear in the rules concluding about C, rules about G are similarly examined. Rules about Z and P are not important at this

point since they appear later in the alphabet than H and will be tested later. If later tests find that one of them uses H in their antecedent list then that rule will be eliminated. The decision to eliminate the current rule (Z C P G ==> H) rather than the previous rule (X H S D ==> C) was made arbitrarily, recognizing the fact that the previous rules had already been check for cycling and were all right.

Control over the number of rules in the knowledge base was lost by this method because many rules were being eliminated. In some cases, an entire ruleset (for example, all rules concluding about H) was removed. This means that either H would have to be tagged undeterminable, or would have to be given a value of true or false since there were no rules to conclude about it.

The method chosen for rule generation allowed second and higher order cycling to occur by keeping the antecedent bucket the same throughout the rule generation. This method differs from the previous one in that the check for first order cycling was done as the antecedents were generated.

For example, suppose the antecedent candidate C is generated for the current rule concluding about H. To test for first order cycling, the antecedents used in the rules concluding about C are scanned for an H. If an H is used as an antecedent in a rule concluding about C, then C is discarded as a candidate for an antecedent in the current rule being generated. If the antecedent candidate generated had been an I, or any letter later in the alphabet than H,

it would have been accepted as an antecedent. The reason for this is because the rules concluding about I - Z have not been generated yet and do not pose a threat to the cycling problem at this point.

b. <u>Second Order Cycling</u> is also expected. An example of such is:

B C D ==> A E F G ==> B A H J ==> E.

Since E does not appear in the antecedent list for A, this set of rules pass the first order cycling test. The first two rules pass the second order test, but the third rule causes a cycle. For example, B is needed to conclude A, E is needed for B, but A is needed for E. It is possible that a cycle would not occur until the fourth, fifth, sixth or higher rule was generated.

It was decided that the amount of time spent checking for cycles higher than first order would be greater than the benefits gained from such a check. Instead, the inference engine was given the responsibility to check for higher order cycling and to take note when this occurred.

3. Random Number Generator

Unfortunately, IQLISP does not have a built in random number generator. However, it does have a function called DTIME that generates hundredths of seconds since midnight, which would be something like 58245 at 9:00 a.m.. The function DTIME produced rules like (E F G H ==> A) which does not give the appearance of being randomly generated. The reason for this is due to the fact that the function DTIME is linear. Adding a delay between generation times eliminated adjacent antecedents but was not effective in altering the linearity.

The next random number generator examined was one that is currently called RANDU and is as follows:

SUBROUTINE RANDU (IX, IY, YFL)

IY = IX * 65539 IF (IY) 5,6,6 5 IY = IY + 2147483647 + 1 6 YFL = IY YFL = YFL * .4656613E-9 RETURN END

It involves very large numbers and relies on the fixed overflow mechanisms of the system. This just means that it is possible to get negative numbers which is the reason for the check for a negative number in line 3. However, IQLISP will allow a number to have 77000 digits and will most likely run out of working memory before a number is generated. An error such as "stack exceeds 64 k " will appear on the screen when this happens.

Rather than trying to modify this algorithm, another was chosen and is as follows:

```
SUBROUTINE RANDOM (A)

MULT = 25211

BASE = 32768

A = MOD [ (A * MULT) , BASE ]

RETURN

END
```

The % function is the IQLISP equivalent to the MOD function which takes the remainder of a division.

This method generates a number between 0 and 1, therefore requiring the FLOAT.LSP package to be loaded into memory before running. The random number is then multiplied by one plus the interval of numbers desired, and the integer portion is added lower bound giving an antecedent candidate, for example, between A and Z.

4. Parameters

As mentioned in the section on rule building, a user has a certain amount of control over the number of conclusions and the number of rules per conclusion wanted. For example, suppose rules to conclude about (A - J) and four rules for each conclusion are desired. This way, each conclusion (A - J) will have four rules concluding about it.

The number of antecedents per rule was arbitrarily chosen randomly to be between two and four. The random number generator RANDOM, mentioned previously, is used to produce this number.

After the antecedent has passed all of the tests,

another number between one and ten is randomly generated. The purpose of this number is to decide whether to add a "not" symbol to the antecedent just generated. If the number generated is one, then the antecedent is negated. Therefore, on the average, every tenth antecedent generated will be transformed into its negation. This number may be increased or decreased if other distributions are desired.

The symbol chosen to represent a "not" is the pound sign (#). An ideal symbol, the tilde ~, was initially chosen, but IQLISP has its own system meanings for the tilde. It uses it to continue a line and is recognized as a comment delimiter.

B. RULE COMPILER

As mentioned in the literature review, it is possible to incorporate much of the search strategy at compile time. The goal is to organize the rules in a way that decreases the amount of work the inference engine has to do at consultation time.

The amount of knowledge incorporated into the compilation depends on the heuristic being used. The heuristic's responsibility is to determine the order in which the antecedents will be tested. The heuristics will be discussed in further detail in the next section. However, in this section, the selection of an antecedent is assumed to be arbitrary. 1. Algorithm

The basic rule compiler algorithm is :

Compile (ruleset)

For all R in ruleset for which the list of antecedents is empty, do

- (1) output conclusion of R;
- (2) remove R from ruleset.

While R is non-empty do

- (3) select an antecedent, say C, from the list of antecedents of some rule;
- (4) output a branch, using antecedent C as the conditional;
- (5) on the true side of the branch; compile all rules that contain an antecedent C after deleting C from each;
- (6) on the negated side of the branch;
 compile all rules that contain an antecedent
 #C after deleting #C from each;
- (7) remove from ruleset those rules compiled in(5) and (6).

A ruleset is a set of rules all concluding about the same thing. For example, a ruleset concluding about A might be as follows:

[((H V I d)A) ((E H)A) ((C D #H)A) ((B C E F)A)] The steps to compile this ruleset are given below The ruleset is not empty, so steps 1 and 2 are skipped.

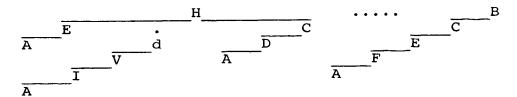
- 3. Suppose H is chosen.
- 4. H is the branch output (it is actually the top of the tree).
- 5. On the true side, compile the new ruleset which is: [((V I d) A) ((E) A)].

This ruleset is not empty, skip steps 1 and 2.

- 3. Suppose E is chosen.
- 4. A branch for E is output.
- 5. On the true side of E, compile new ruleset:
 [(() A)].
 - 1. The new ruleset is empty, so the conclusion A is output.
- 6. There is no negated side (rules involving #E).

7. New ruleset : [((V I d) A)].

The process continues, putting out branches for the rest of the antecedents in the ruleset, until the tree looks like:



where the indicates a sequential list of trees.

2. Data Structures

a. <u>Decision Trees</u> are represened as follows : (antecedent (if ant. is true) (if ant. is false)). A conclusion may have more than one decision tree associated with it. In the previous example, there were two associated with A - one tree with H at the top, and one with B at the top. Therefore, each conclusion is associated with a decision tree list.

b. <u>Decision Tree Lists</u> are bound to the name of the conclusion which is represented by an asterisk with the consequent. For example, *A = [(H (if H is true) (if H is false)) (B (if B is true) (if B is false))]. To find out if H is true or false, its decision tree has to be examined which may itself be a decision tree list.

This structure eliminates all negation symbols from the rulesets by having a true decision tree list and a false decision tree list. The need for sequential trees was discovered as more complicated rulesets were experimented with, and this structure will handle any number of sequential trees.

C. HEURISTICS FOR ANTECEDENT SELECTION

The most important part of the compile algorithm is the selection of the antecedent to put out as a branch. In the section on compiling, the selection of an antecedent was assumed to be arbitrary. This section presents two heuristics that were investigated, each of which approach the selection process in very different ways.

The first heuristic was used by van Melle in EMYCIN. It is straightforward and uses no knowledge about the relationship among the rules. The second heuristic incorporates knowledge about the rules generated, specifically, the average number of antecedents required to conclude about a parameter.

1. Most Often Occurring

The reasoning behind this heuristic is that an antecedent appearing in the most rules must be important and therefore should be placed at the top of the tree. To exhibit this heuristic, assume the same ruleset as before:

[((H V I d)A) ((E H)A) ((C D #H)A) ((B C E F)A)]. A new list is made from the ruleset list containing each antecedent and the number of times it appears in the ruleset. For example,

[(H 3) (V 1) (I 1) (d 1) (E 2) (C 2) (D 1) (F 1)] Note that the interest is in how many times an antecedent is used, disregarding how it was used (negation). Therefore, H would be chosen as the first antecedent (branch).

2. Minimum Average Antecedent

This heuristic examines the relationship between the antecedent in the current ruleset with those in the rest of the knowledge base.

The reasoning is to choose the antecedent that, on the average, requires the least amount of work to determine its value. For example, if the average number of antecedents needed to conclude G is two and the average number of antecedents needed to conclude H is four, then G will be chosen.

The number of antecedents required for a particular

rule is readily available from the array ANT_INFO as discussed in the section on rules. The total number of antecedents minus the number of knowns is the number of antecedents that will be inferred to conclude a particular rule. For example, in the rule ((H V I d) X) the total number of antecedents is four, the number of knowns is one. Therefore, the number of antecedents needed to conclude X is three.

An extremely misleading assumption to make is that the number of antecedents required to conclude X is only three. This implies that the number of antecedents required to conclude H is one, the number to conclude V is one and the number to conclude I is one which is possible, but unlikely.

To rectify this problem, the rules with H, V, and I as consequents must be examined. Then the rules whose consequents are the antecedents in H, V, and I must be examined and new averages calculated. As can be seen, this is an n-order problem. For this research, averages were calculated three times in addition to the averages calculated during generation of the rules. The decision to calculate averages an additional <u>three</u> times was chosen arbitrarily. However, consultation results (number of IFs) were very close, in some cases identical, between the knowledge bases compiled with second and third averages indicating a point of diminishing return.

D. INFERENCE ENGINE

The inference engine's job, if the compiler performed correctly, is trivial. It retrieves the decision trees built by the compiler, examining either the true side or false side of the tree depending on the value of the current antecedent. Recall that a decision tree takes the form: *A = ((a (true side) (false side)) (sequential trees)), where a = antecedent, (true side) or (false side) can be a list of decision trees. So if the antecedent is true, the inference engine recurs with the CAR of the true side. If that returns false or undeterminable, the next tree (still on the true side of the antecedent) is examined. If there are no more trees (possible ways to conclude the antecedent) then "undeterminable" is returned from the true side and also to *A if there are no more sequential trees. Therefore, using the example in section B.l., the compiled decision trees are structured as follows: *A = [(decision tree #1) (decision tree #2)]

- = [(H (true side)(false side)) (B (true side) (false side))]
 - = [(H ((decision tree 1) (decision tree 2)) (false
 side))
 (B (true side) (false side = nil))]
 = [(H ((E (A) nil)(d ((V ((I (A) nil)) nil)) nil))
 - (C ((D (A) nil)) nil) <= false side of H
) <== end of decision tree # l
 (B ((C ((E ((F (A) nil)) nil)) nil)) nil)]</pre>

IV. EXAMPLE CONSULTATION

Suppose the following set of rules are under consideration :

[(((b c d)F) ((I H)F)) (((F #H) G) ((a J)G))]

(((e J)H) ((g I b)H)) (((#G a)I) ((#f e b)I))

(((F e)J) ((#F I)J))].

Note that each ruleset has two rules concluding about the conclusions F, G, H, I, and J. They will be compiled under the assumption that antecedents will be chosen in the order of (F #F G #G H #H I #I J #J). However, if any rule contains a known (lower case letter) it will be chosen first.

The consultation begins with the user specifying the values of the lower case antecedents. For this research, all lower case letters were set to be true. This was to enable consistency since different combinations of "trues" and "falses" will produce different results. For this example, it is assumed that the goals for the consultation are G and H.

*J = (((e ((F (J) NIL)) NIL) (F NIL ((I (J) NIL))))))

INFERENCE ENGINE LOG: (SET TREES '(a b c d e f g) '**TRUE**) (INFER (QUOTE G)) a is known to be true Attempting to satisfy J e is known to be true Attempting to satisfy F b is known to be true c is known to be true d is known to be true F is deduced to be true J is deduced to be true G has been deduced to be true after 7 IFs (INFER (QUOTE H)) b is known to be true g is known to be true Attempting to satisfy I a is known to be true Attempting to satisfy G a is known to be true Attempting to satisfy J e is known to be true Attempting to satisfy F b is known to be true c is known to be true d is known to be true F is deduced to be true J is deduced to be true G is deduced to be true I cannot be determined b is known to be true e is known to be true f is known to be true I cannot be determined I cannot be determined I cannot be determined H cannot be determined H cannot be determined e is known to be true J is known to be true

H has been deduced to be true after 16 IFs

V. SUMMARY OF RESULTS AND CONCLUSIONS

This paper presented a system, CIEGEN, developed as a tool for building expert systems on a microcomputer and as a valuable aid for research on rule compilation. It was written in IQLISP and implemented on the IBM PC.

The main purpose of CIEGEN's rule generator is for use in research on rule compilation because it allows the generation of parameterized knowlege bases.

In CIEGEN, the knowledge base can be described by the number of parameters to be concluded about, the number of knowns per knowledge base, the number of rules per conclusion, and the number of antecedents per rule. These characteristics are input to the Rule Builder which keeps track of the number of knowns per rule, the total number of antecedents per rule, the unknowns or antecedents to be inferred per rule, and the average number of antecedents used per rule as rules are generated. In order for a generated antecedent to be accepted, it must guarantee not to cause first order cycling. If first order cycling is detected, it is discarded and another antecedent is generated.

After the rules have been generated, they are compiled into a form that will hopefully execute more efficiently. Rule compilation is the process of transforming a knowledge base consisting of rules to a knowledge base of decision trees. The transformation effectively takes place by choosing an antecedent to be placed at the top of the tree,

gathering all of the rules using this antecedent, choosing another antecedent as a branch, gathering those rules using this antecedent, continuing the process until a rule has had all of its antecedents chosen as branches, then printing the conclusion as a leaf node. At this point, those rules containing the negation of the antecedents output as branches are gathered and the process is repeated, forming the right side of the tree. As mentioned previously, this technique has been implemented in EMYCIN, tested against the technique of interpreting rules, and shown to cut interquestion time close to half for the expert systems MYCIN, PUFF, and SACON.

The efficiency of the resulting decision trees depends on the method of selecting the antecedents. The author developed a heuristic (Minimum Average Antecedent) for selecting antecedents based on knowledge gathered at generation time. The heuristic forces those antecedents requiring the least amount of work, on the average, to be chosen first. The amount of work is determined by the number of antecedents needed to be inferred in order to conclude the antecedent in the current rule.

This heuristic was compared with a heuristic (Most Often Occurring) used by EMYCIN. The Most Often Occurring heuristic chooses the antecedent that appears in the most rules. The heuristic is based upon the idea that if it is used in the most rules, it must be important and should be placed nearest the top of the tree.

The efficiency of the two heuristics was measured by the number of IFs executed during a consultation with each knowledge base. The performance of a consultation is the responsibility of CIEGEN's inference engine. The inference engine consists of a series of procedures which examine the decision trees built by the rule compiler and counts the number of antecedents it has to infer before a conclusion can be made. The inference engine retrieves the decision tree list of the parameter typed in by the user (what is to be inferred, which can be the negation of a parameter), then recursively examines the decision trees of the antecedents in the original tree until a conclusion is made.

An antecedent's decision tree will only be examined once and at that time the value is bound to the antecedent. It is possible to need the value of an antecedent that is being inferred, thus creating a cycle. When this happens, the inference engine notes that the current path cannot be continued and retrieves the next decision tree in the list if one exists.

The inference engine traces the paths led by the decision trees and outputs the total number of nodes it visited with the value of the parameter being inferred. The value of the parameter will either be "true" or "undeterminable". The parameter will be undeterminable if its decision trees were involved in cycles preempting it from making a conclusion, or if the only path available required the parameter to be false.

The fact that the parameters are never concluded to be false (unless the negation of the parameter is being inferred) is not obvious with the emphasis on a generic knowledge base. The parameters could represent parts of more complex rules. For example, concluding "A" to be true could represent the conclusion "it is true that the bacteria is not present." Similarly, concluding "B" to be true could represent that the bacteria is present. A rule requiring that the bacteria is not present would use "A" instead of "#B" and vice versa. Therefore, it is possible to generate a realistic knowledge base with all non-negated antecedents and conclusions that are true.

Some parameters describing the knowledge base were held constant and some were allowed to vary. For each knowledge base, four rules were generated for each letter in the alphabet (A-Z) and two to four antecedents were created for each rule. These numbers were arbitrarily chosen, as a knowledge base consisting of 104 rules was considered to be a reasonable size. Knowledge bases permitting negations were generated with the number of knowns varying among 5, 7, and 14. Another set of knowledge bases were generated witholding negations and varying the number of knowns among 5, 7, 14, and 21.

For each number of knowns, there were five knowledge bases generated and tested. The intent for varying the number of knowns was to test whether either heuristic would be affected by giving it more information. The goals of

each consultation were assumed to be A-E, so the averages, shown in Table I, are based on the total number of IFs required to infer these five parameters. The right side of Table I shows the results of consultations performed on knowledge bases without negations in the rules. The reason for restricting negations from the knowledge base was because the second heuristic Minimum Average Antecedent did not take into account the fact that parameters are never deduced to be false.

Overall, the Most Often Occurring heuristic performed consultations in the least number of IFs. As the number of knowns increased, Most Often Occurring's average increased. The author suggests that this is because there are fewer common antecedents as the number of knowns are increased. In other words, the number of antecedents to choose from ranged from 26 to 28 to 35 which means that there will be less repetitive antecedents in the rules. However, in the case where the number of knowns and the number of parameters to be inferred were equal, Most Often Occurring's average dropped rather than increased. This may be due to the fact that because there are so many knowns in the knowledge base the actual cause of the decrease is obscured.

The Minimum Average Antecedent heuristic explicitly put knowns closest to the top of the decision tree. This means that if there are knowns in a rule, they will be tested first. Therefore, it is expected that as the number of knowns increases, the overall average will decrease. The

TABLE I

CONSULTATION RESULTS

AVERAGE NUMBER OF IFS

	Antecedent selection interval (negations permitted)			Antecedent selection interval (negations withheld)			
HEURISTIC	26	28	35	26	28	35	42
Most Often Occurring	75	81	84	76	118	123	65
Minimum Average Antecedent	152	132	98	123	126	91	74

Often Occurring heuristic did not explicitly place knowns first, it merely placed rules containing knowns closest to the top.

As mentioned previously, the Minimum Average Antecedent algorithm did not take into account the possibility of failures. It makes the assumption that the conclusions made will be true. Since parameters are never concluded to be false, if a rule requires the negation of an antecedent to be true, the rule will fail. In this case the average number of antecedents to conclude a parameter no longer makes sense. Possibly incorporated into this heuristic should be a penalty for those parameters that have "nots" present with their antecedents. The other alternative is the one previously mentioned, letting the parameters or antecedents represent things more complicated than true or false.

Data concerning times of the system were also collected. The average amount of time to generate a knowledge base was twenty minutes. The average amount of time to compile a knowledge base was 31 minutes for the Most Often Occurring heuristic and 29 minutes for the Minimum Average Antecedent. The slight difference may be attributed to the calculations that are done by the Most Often Occurring heuristic tallying the number of times an antecedent appears in the rules. An additional thirteen minutes were required to calculate averages for the Minimum Average Antecedent heuristic. An interesting note to make is that a slightly more complicated heuristic took almost 40% more time to compile.

As demonstrated by the comparison of these two heuristics, CIEGEN is a convenient tool for research and an in depth analysis of heuristics. Information about the rules such as the number of antecedents, values of antecedents, number of rules, and number of knowns is readily available and can easily be varied or held constant. The system is not only domain independant, but is heuristic independant so that heuristics can be easily inserted into the system for testing and comparison with other compilation heuristics.

31

VI. FURTHER RESEARCH SUGGESTIONS

CIEGEN provides a base for interesting work in at least three areas. A natural extension to the work presented in this thesis are the following areas:

As a research tool ~

- a statistical analysis of the interactions between parameters
- a study of additional heuristics
- an addition to the Minimum Average Antecedent heuristic to incorporate negations
- a study of the complexity of heuristics vs. compile
 time

As groundwork for an expert system -

- explanation capabilities for the inference engine
- the incorporation of uncertainty in the rules
- the allowance of more complex rules with multiple conclusions

In the generation of rules -

- the incorporation of more knowledge about cycling, rather than just checking for first order
 - the incorporation of a learning mechanism

CIEGEN provides the groundwork for each of these areas.

BIBLIOGRAPHY

- 1. Barr, Avron and Edward Feigenbaum. <u>The Handbook of of Artificial Intelligence</u>, vol. 1. California: William Kaufmann, Inc, 1981.
- Hendrix, Gary G. and Earl D. Sacerdoti, "Natural Language Processing: The Field in Perspective," <u>Byte</u> 6, 9 (1981), 304-352.
- 3. Winograd, Terry. Language as a Cognitive Process. Massachusetts : Addison-Wesley Publishing Co., 1983.
- 4. Prendergast, Dan. "A General Purpose Robot Control Language," Byte, 9, 1 (1984), 122-133.
- 5. Cohen, Paul R. and Edward Feigenbaum. <u>The Handbook of</u> <u>Artificial Intelligence</u>, vol. 3. California: William Kaufman, Inc., 1982.
- 6. Feigenbaum, E. A. "The Art of Artificial Intelligence : Themes adn Case Studies of Knowledge Engineering," Report : STAN-CS-77-621, Computer Science Department, Stanford University, Stanford, California.
- 7. Barr, Avron and Edward Feigenbaum. <u>The Handbook of</u> <u>Artificial Intelligence</u>, vol.2. California: William Kaufman, Inc., 1982.
- 8. Buchanan, Bruce and Edward Feigenbaum. "DENDRAL and META-DENDRAL: Their Application Dimension," <u>Artificial</u> Intelligence, 11 (1978), 5-24.
- 9. Buchanan, Bruce and Tom M. Mitchell. "Model Directed Learning of Production Rules." Report: HPP 77-6. Heuristic Programming Project, Computer Science Department, Stanford University, Stanford, California.
- 10. Hayes-Roth, Frederick and others. <u>Building Expert</u> <u>Systems</u>. <u>Massachusetts</u>: Addison-Wesley Publishing Co., 1983.
- 11. Farley, A. M. "Issues in Knowledge Based Problem Solving," <u>IEEE Transactions System Man and Cybernetics</u> SMC-10, 8 (1980), 446-459.
- 12. Van Melle, W. "MYCIN : A Knowledge Base Consultation Program for Infectious Disease Diagnosis," <u>International</u> <u>Journal of Man-Machine Studies</u>, 10 (1978), 313-322.

- 13. Van Melle, W. "A Domain Independant Production Rule System for Consultation Programs," STAN-CS-80-820, Stanford, California, June, 1980.
- 14. Forgy, Charles, L. "Rete: A Fast Algorithm for the Many Pattern / Many Object Pattern Match Problem," <u>Artificial</u> <u>Intelligence</u>, 19 (1982), 17-37.
- 15. Davis, Randall. "Interactive Transfer of Expertise: Acquisition of New Inference Rules," <u>Artificial</u> Intelligence, 12, 2 (1979), 121-157.
- 16. Davis, Randall and Douglas Lenat. <u>Knowledge Based</u> <u>Systems in Artificial Intelligence</u>. New York: <u>McGraw-Hill International Book Co.</u>, 1982.
- 17. <u>IQLISP</u> <u>Reference</u> <u>Manual</u>. Washington: Integral Quality, 1983.

VITA

Jayne Denise Ward was born on February 16, 1960 in Springfield, Missouri where she also received her primary and secondary education. She graduated cum laude from Southwest Missouri State University in Springfield, Missouri with a Bachelor of Science in Computer Science and Mathematics in May, 1982.

She has been enrolled in the Graduate School of the University of Missouri-Rolla since August, 1982. Since June, 1983 she has been employed by the United States Geological Survey in Rolla, Missouri and has held a Graduate Teaching Assistantship in the Department of Computer Science from August, 1983 to May, 1984. During May and June, 1984, she has held a Research Assistantship in the Department of Computer Science. She is a member of the honor societies Upsilon Pi Epsilon, Kappa Mu Epsilon, and Sigma Pi Sigma as well as a student member of the American Association for Artificial Intelligene and the Association for Computing Machinery.

35

APPENDIX A

CIEGEN PACKAGES

PACKAGE "COMPILER.LSP"

RC

Arguments : PARTITION LIST -- a list of partitions, FILENAME -- filename (in double quotes) specifying a disk file to receive the listing produced by the compiler. the user Called by : Calls DO : RETURNS : the compilation. RC is the function that causes rules to be compiled. The user should type, for example, (RC PL "B:RC3") which will cause the partition list PL to be compiled and output to be sent to the screen and to the file B:RC3 (DEF 'RC [LAMBDA (PARTITION LIST FILENAME) [PROG () (SETQ *PARMS* NIL) (SETQ SESSIONFILE (OUTPUT FILENAME)) (PRINTC '"RULES SUBMITTED:" SESSIONFILE) (PRINTC '"RULES SUBMITTED:") (TERPRI NIL SESSIONFILE) (TERPRI NIL SESSIONFILE) (TERPRI) (TERPRI) (PP PARTITION LIST SESSIONFILE) (PP PARTITION LIST) (TERPRI NIL SESSIONFILE) (TERPRI) (PRINTC '"RESULTS OF COMPILATION:" SESSIONFILE) (PRINTC '"RESULTS OF COMPILATION:") (TERPRI NIL SESSIONFILE) (TERPRI) (DO PARTITION LIST) (CLOSE SESSIONFILE) (TERPRI) (RETURN 'END OF COMPILATION)]]) DO PARTITION LIST -- a list of rule partitions Arguments : (a complete set of input rules grouped by conclusion parameters) Called by : RC, itself : COMPILE, itself Calls passes each partition to COMPILE. DO puts an RETURNS : extra set of parentheses around the output of COMPILE making it a decision tree list

```
(DEF 'DO
 '[LAMBDA (PARTITION_LIST)
   [PROG (P)
       (COND
           [(NULL PARTITION LIST)
             (RETURN 'DONE)]
           Гт
             (SETQ P
                 (CAR PARTITION LIST))
             (SETQ OUTFILE
                 (OUTPUT "B:OUTFILE"))
             (PRINC '"(" OUTFILE)
             (COMPILE P)
             (PRINC '")" OUTFILE)
             (CLOSE OUTFILE)
             (SETQ INP
                 (INPUT "B:OUTFILE"))
             (SETQ TREENAME
                 (READLIST (CONS '"*"
                                (EXPLODE (CADAR P)))))
             (SET TREENAME
                 (READ INP))
             (CLOSE INP)
             (PRINT TREENAME SESSIONFILE)
             (PP (EVAL TREENAME) SESSIONFILE)
             (PRINT TREENAME)
             (PP (EVAL TREENAME))
             (SET (READLIST (CONS '"!"
                                 (EXPLODE (CADAR P))))
                 (EVAL TREENAME))
             (SETQ *PARMS*
                 (CONS (CADAR P) *PARMS*))
             (DO (CDR PARTITION LIST))
             (RETURN 'DONE)])]]
                    COMPILE
Arguments :
              RULESET -- A rule partition (a list of rules
                         concluding about a particular ant)
Called by :
              DO, itself
              CONCLUDE ALL W SATISFIED PREMISES, ALL NOT
Calls
           :
              SATISFIED, SELECT A CLAUSE, UNNEGATED, DEL
              RULE SET, FIND ALL, REMOVE, NEGATION
```

RETURNS : a decision tree list

(DEF 'COMPILE

'[LAMBDA (RULESET)

[PROG (C K)

(CONCLUDE_ALL_W_SATISFIED_PREMISES RULESET) (SETQ K (ALL_NOT_SATISFIED RULESET)) CONTINUE (COND

```
[(NULL K)
  (RETURN 'DONE)]
[Т]
  (SETQ C
  (SELECT_A_CLAUSE_K))
(PRINC '"["OUTFILE)
  (COND
      [(UNNEGATED C)
         (PRIN C OUTFILE)
(PRINC '"(" OUTFILE)
         (COMPILE (DEL RULESET C
                     (FIND ALL C K)))
         (PRINC '")" OUTFILE)
         (SETQ K
             (REMOVE C K))
         (PRINC '"(" OUTFILE)
         (COMPILE (DEL_RULESET (NEGATION C)
                       (FIND ALL (NEGATION C) K)))
         (PRINC '")" OUTFIL\overline{E})
         (SETQ K
             (REMOVE (NEGATION C) K))]
      [Т]
         (PRIN (NEGATION C) OUTFILE)
         (PRINC '"(" OUTFILE)
         (COMPILE (DEL RULESET (NEGATION C)
                       (FIND ALL (NEGATION C) K))
         (PRINC '')' OUTFILE)
         (SETQ K
         (REMOVE (NEGATION C) K))
(PRINC '"(" OUTFILE)
         (COMPILE (DEL RULESET C
                      (FIND ALL C K))
         (PRINC '")" OUTFILE)
         (SETQ K
            (REMOVE C K))])
  (PRINC '"]" OUTFILE)
  (GO CONTINUE)])])
```

NEGATION

Called by Calls		ANT an antecedent COMPILE system functions only the negation of the antecedent. This function is used by the compiler to specify rules which contain the negation of a particular ant.
(DEF 'NEGA '[lambda (
(COND	1 1	• /
[(NU	JLL	ANT)
	[L]	
[(EC	(UA	L (CAR (EXPLODE ANT))

· " # ") (READLIST (CDR (EXPLODE ANT)))] Гт (READLIST (CONS '"#" (EXPLODE ANT)))])]) ALL W SATISFIED PREMISES RULESET -- a rule partition Arguments : Called by : CONCLUDE ALL W SATISFIED PREMISES Calls : itself RETURNS a list of all conclusions which have satis-: fied premises (a nil antecedent list). (DEF 'ALL W SATISFIED PREMISES '[LAMBDA (RULESET) (COND [(NULL RULESET) NIL] [(NULL (CAAR RULESET)) (CONS (CADAR RULESET) (ALL W SATISFIED PREMISES (CDR RULESET)))] Гт (ALL W SATISFIED PREMISES (CDR RULESET))])]) CONCLUDE ALL W SATISFIED PREMISES Arguments : RULESET -- a rule partition Called by : COMPILE : ALL W SATISFIED PREMISES, PRINT EACH Calls : construction of leaf nodes. Prints each RETURNS element returned by ALL W SATISFIED PREMISES (DEF 'CONCLUDE ALL W SATISFIED PREMISES '[LAMBDA (RULESET) (COND [(NULL (ALL W SATISFIED PREMISES RULESET)) NIL] [T] (PRINT EACH (ALL_W_SATISFIED_PREMISES RULESET) OUTFILE)] JJJ ALL NOT SATISFIED Arguments : RULESET -- a rule partition Called by : COMPILE, itself Calls itself : : everything in the rule partition that does not RETURNS

(DEF 'ALL NOT SATISFIED '[LAMBDA (RULESET) (COND [(NULL RULESET) NIL] [(NULL (CAAR RULESET)) (ALL NOT SATISFIED (CDR RULESET))] Гт (CONS (CAR RULESET) (ALL NOT SATISFIED (CDR RULESET)))])]) REMOVE Arguments : CLAUSE -- an antecedent, RULESET -- a rule partition Called by : COMPILE, itself : ANT IS IN RULE, REMOVE Calls : a partition list which is RULESET without the RETURNS rules that reference CLAUSE (DEF 'REMOVE [LAMBDA (CLAUSE RULESET) (COND [(NULL RULESET) NIL] [(ANT IS IN RULE CLAUSE $\overline{(CAR RULESET)}$ (REMOVE CLAUSE (CDR RULESET))] ГΤ (CONS (CAR RULESET) (REMOVE CLAUSE (CDR RULESET)))])]) PRINT EACH Arguments : LST -- a list of parameters Called by : CONCLUDE ALL W SATISFIED PREMISES, itself : itself Calls : prints each of the elements in LST to OUTFILE RETURNS (DEF 'PRINT EACH '[LAMBDA (LST) (COND [(NULL LST) 'DONE] [т] (PRIN (CAR LST) OUTFILE) (PRINT EACH (CDR LST))])]) UNNEGATED

Arguments : ANT -- an antecedent

```
Called by : COMPILE
Calls : LISP builtin functions only
RETURNS : T if the antecedent is unnegated, NIL
otherwise
(DEF 'UNNEGATED
'[LAMBDA (ANT)
 (COND
    [(EQUAL '"#"
        (CAR (EXPLODE ANT)))
        NIL]
    [T
        T])])
```

INFER

```
PARM -- a parameter
Arguments :
Called by :
               The user.
               UNNEGATED, IS KNOWN TO BE TRUE, IS KNOWN TO
Calls
           Ξ
               BE FALSE, IS KNOWN TO BE UNDET, INFER LIST,
               END CONSULTATION, OPPOSITE, NEGATION
               INFER is the top level function for the
RETURNS
           :
               inference engine. The user calls the function
               with the parameter whose truth value is
               desired. The inference engine then displays
               a trace of the traversal of the knowledge base
               as it infers the parameter.
(DEF 'INFER
'[LAMBDA (PARM)
  [PROG (TEMP)
       (TERPRI NIL IELOG)
       (SETQ IM_WORKING ON NIL)
       (COND
           [(UNNEGATED PARM)
              (COND
                  [(IS KNOWN TO BE TRUE PARM)
                    (RETURN (END CONSULTATION PARM
(T 1)))]
                  [(IS KNOWN TO BE FALSE PARM)
                     (RETURN (END CONSULTATION PARM
                                  <sup>+</sup>(F 1)))]
                  [(IS KNOWN TO BE UNDET PARM)
                     (RETURN (END CONSULTATION PARM
(U 1)))]
                  [Т]
                     (SETO TEMP
                         (INFER LIST PARM
                              (EVAL (READLIST (CONS '"*"
                                              (EXPLODE PARM))))
                         0 0 NIL))
                     (RETURN (END CONSULTATION PARM TEMP))])]
           Ет
              (COND
                  [(IS KNOWN TO BE TRUE (NEGATION PARM))
                     (RETURN (END CONSULTATION PARM
(F 1)))]
                  [(IS KNOWN TO BE FALSE (NEGATION PARM))
                     (\overrightarrow{RETURN} (\overrightarrow{END} \overrightarrow{CONSULTATION} \overrightarrow{PARM} (T 1)))]
                  [(IS KNOWN TO BE UNDET (NEGATION PARM))
                     (RETURN (END CONSULTATION PARM (U 1)))]
                  [Т]
                    (SETQ TEMP
```

INFER LIST

PARM -- a parameter, DT LIST -- a decision Argruments : tree list which is to be used to determine the truth of the parameter, COUNT -- the count of IFs processed so far, SPACE COUNT -controls the indentation of messages printed as the inference process progresses, GIVE UPcontrols whether a message should be printed indicating that the parameter cannot be determined. It is possible to arrive at a place in the knowledge base where the DT LIST NIL, but there are still decision trees to be searched. (This is due to having sequences decision trees.) GIVE UP will be T if INFER LIST has just called itself directly and will be NIL otherwise (INFER2 made the call). INFER, INFER2, itself Called by : PRINT LINE, INFER2, itself Calls : if the parameter is determined to be true or RETURNS : false, then the search stops and INFER LIST returns the truth and count value fount. If the parameter cannot be determined using one decision tree, the remaining decision trees used until the parameter is determined or there are no more decision trees in the list. (DEF 'INFER LIST [LAMBDA (PARM DT LIST COUNT SPACE_COUNT GIVE_UP) [PROG (TEMP) (RETURN (COND [(NULL DT LIST) (COND [GIVE UP (PRINT LINE SPACE COUNT PARM '"cannot be determined") (SET (READLIST (CONS '"*" (EXPLODE PARM))) '**UNDET**) (SETQ IM WORKING ON (DEL CURR GOAL IM WORKING ON PARM)) (CONS U (LIST (ADD1 COUNT)))] Гт (CONS 'U (LIST COUNT))])]

[(EQUAL 'U (CAR (SETQ TEMP (INFER2 PARM (CAR DT LIST) COUNT SPACE_COUNT)))) (INFER_LIST PARM (CDR DT_LIST) (CADR TEMP) SPACE_COUNT T)] [T TEMP]))]])

INFER2

Arguments : PARM -- the parameter to be determined, DT -a decision tree to be used in determining PARM COUNT -- the count of IFs seen so far, SPACE COUNT -- indentation value for messages Called by : INFER LIST IS KNOWN TO BE TRUE, IS KNOWN TO BE FALSE, Calls : IS KNOWN TO BE UNDET, PRINT LINE, INFER LIST the truth and count pair indicating the result RETURNS : of its search. (see INFER_LIST) It searches a single decision tree. (DEF 'INFER2 '[LAMBDA (PARM DT COUNT SPACE COUNT) [PROG (TEMP) (RETURN (COND [(ATOM DT) (COND [(EQUAL PARM DT) (SET (READLIST (CONS '"*" (EXPLODE PARM))) '**TRUE**) (CONS 'T (LIST COUNT))] [T] (TERPRI) (SPACES SPACE COUNT) (PRINC '"malformed decision tree") (COND [ECHO ON? (TERPRI NIL IELOG) (SPACES SPACE COUNT IELOG) (PRINC ' "malformed decision tree") IELOG)] [т] NIL]) (LIST 'U COUNT)])] [(IS KNOWN TO BE TRUE (CAR DT))

```
(PRINT LINE SPACE COUNT
      (C\overline{A}R DT)
       '"is known to be true")
  (INFER LIST PARM
      (CADR DT)
      (ADD1 COUNT) SPACE COUNT NIL)]
[(IS KNOWN TO BE FALSE (CAR DT))
  (PRINT LINE SPACE COUNT
      (C\overline{A}R DT)
       '"is known to be false")
  (INFER LIST PARM
      (CADDR DT)
      (ADD1 COUNT) SPACE COUNT NIL)]
[(IS_KNOWN_TO_BE_UNDET (CAR DT))
  (PRINT LINE SPACE COUNT
      (CAR DT)
       '"is known to be undeterminable")
  (CONS 'U
      (LIST (ADD1 COUNT)))]
  (COND
      [(NULL (MEMBERS (CAR DT)
                          IM WORKING ON))
         (SETQ IM WORKING ON
            (CONS (CAR DT IM WORKING ON))
         (TERPRI)
         (SPACES SPACE COUNT)
         (PRINC '"Attempting to satisfy ")
         (PRIN (CAR DT))
         (COND
             [ECHO ON?
               (TERPRI NIL IELOG)
               (SPACES SPACE COUNT IELOG)
               (PRINC '"Attempting to
                  satisfy " IELOG)
               (PRIN (CAR DT) IELOG)]
             ГΤ
               NIL])
         (COND
             [(EQUAL 'T
                  (CAR (SETQ TEMP
                      (INFER LIST (CAR DT)
(EVAL (READLIST
                                (CONS '"*"
               (EXPLODE (CAR DT)))) COUNT
                 (+ 2 SPACE COUNT) NIL))))
               (PRINT LINE SPACE COUNT
                    (C\overline{A}R DT)
                '"is deduced to be true")
               (SETQ IM WORKING ON
              (DEL CURR GOAL IM WORKING ON
                        (CAR DT)
               (INFER LIST PARM
                    (C\overline{A}DR DT)
```

[Т]

(ADD1 (CADR TEMP)) SPACE COUNT NIL)] [(EQUAL 'F (CAR TEMP)) (PRINT LINE SPACE COUNT (CAR DT) '"is deduced to be false") (SETQ IM_WORKING_ON (DEL CURR GOAL IM WORKING ON (CAR DT)(INFER LIST PARM (CADDR DT) (ADD1 (CADR TEMP)) SPACE COUNT NIL)] ГТ (LIST 'U (CADR TEMP))])] [T] (PRINT LINE SPACE COUNT (CAR DT) '"is involved in a cycle") (CONS 'U OPPOSITE TRUTH AND COUNT -- a pair, or list of two Arguments : elements, containing the "truth" of a parameter and the count of the number of IFs processed in extablishing the truth. Called by : INFER Calls builtin functions only : RETURNS the negation of the parameter and the same : count (DEF 'OPPOSITE [LAMBDA (TRUTH AND COUNT) (COND [(ATOM TRUTH AND COUNT) (PRINT '"malformed truth and count value") (COND [ECHO ON? (PRINT '"malformed truth and count value" IELOG)] [т NIL]) (NIL)] [(EQUAL 'T (CAR TRUTH AND COUNT)) (CONS 'F (CDR TRUTH AND COUNT))] [(EQUAL 'F (CAR TRUTH AND COUNT)) (CONS 'T (CDR TRUTH AND COUNT))] [(EQUAL 'U

```
(CAR TRUTH_AND_COUNT))
TRUTH_AND_COUNT]
[T
(PRINT '"malformed truth and count value")
(COND
       [ECHO_ON?
       (PRINT '"malformed truth and count value" IELOG)]
       [T
            NIL])
(NIL)])])
```

IS KNOWN TO BE TRUE

Called by : IN Calls : bu RETURNS : T th **	RM a parameter FER, INFER2 iltin functions only if the decision tree list corresponding to e parameter PARM is bound to the value TRUE**, corresponding to PARM being known be true. Returns NIL otherwise.
---	---

IS KNOWN TO BE FALSE

(Similar to IS_KNOWN_TO_BE_TRUE)

IS KNOWN TO BE UNDET

(similar to IS_KNOWN_TO_BE_TRUE)

SET TREES

Arguments : PARMLIST -- a list of parameters, VAL -- a value Called by : the user, itself Calls : itself : the decision tree list set equal to the value RETURNS VAL. (DEF 'SET TREES [LAMBDA (PARMLIST VAL) (COND [(NULL PARMLIST) 'FINE] [T] (SET (READLIST (CONS '"*" (EXPLODE (CAR PARMLIST)))) VAL) (SET TREES (CDR PARMLIST) VAL) 'FINE])])

IT

Arguments : none Called by : the user. : TREE INIT Calls : the reinitialization of decision tree lists RETURNS without recompilation of rules. (DEF 'IT '[LAMBDA () (TREE INIT *PARMS*)]) TREE INIT Arguments : PARMLIST -- a list of parameters Called by : IT, itself Calls : itself (DEF 'TREE INIT '[LAMBDA (PARMLIST) (COND [(NULL PARMLIST) 'SURE] ГΤ (SET (READLIST (CONS '"*" (EXPLODE (CAR PARMLIST)))) (EVAL (READLIST (CONS '"!" (EXPLODE (CAR PARMLIST)))))) (TREE INIT (CDR PARMLIST)) 'SURE])])

PRINT_LINE

Arguments : SPACE_COUNT -- how many spaces to indent the

printed line, THING -- the thing to be printed MSG -- a message to be printed with the value of THING Called by : INFER LIST, INFER2, END CONSULTATION Calls builtin functions : RETURNS a line of text to be printed as part of the : trace of execution of the inference engine. (DEF 'PRINT LINE '[LAMBDA (SPACE COUNT THING MSG) [PROG () (TERPRI) (SPACES SPACE COUNT) (PRIN THING) (SPACES 1) (PRINC MSG) (COND [ECHO ON? (TERPRI NIL IELOG) (SPACES SPACE COUNT IELOG) (PRIN THING $I\overline{E}LOG$) (SPACES 1 IELOG) (PRINC MSG IELOG)] [т NIL])]]) END CONSULTATION PARM -- a parameter, PAIR -- a truth and count Arguments : pair INFER Called by : : PRINT LINE Calls prints some final information that appears RETURNS : after the inference engine has accomplished it can. (DEF 'END CONSULTATION [LAMBDA (PARM PAIR) [PROG () (TERPRI) (COND [ECHO ON? (TERPRI NIL IELOG)] $[\mathbf{T}]$ NIL]) (PRINT LINE O PARM (COND [(EQUAL 'T (CAR PAIR)) "has been deduced to be true"] [(EQUAL 'F

(CAR PAIR))

'"has been deduced to be false"]

```
ГΤ
          '"is undeterminable"]))
(PRINC '" after ")
(PRIN (CADR PAIR))
(PRINC '" IFs")
(TERPRI)
(TERPRI)
(COND
    [ECHO ON?
      (PRINC '" after " IELOG)
      (PRIN (CADR PAIR) IELOG)
      (PRINC '" IFs" IELOG)
      (TERPRI NIL IELOG)]
    Гт
      NIL])
(RETURN 'CONSULTATION ENDED)]])
```

NEGATION

Arguments : ANT -- an antecedent Called by : INFER Calls : builtin functions only RETURNS : the negation of the parameter (antecedent). This function allows a user to infer the value of a negated parameter. The unnegated form is inferred and the answer is returned negated. (The decision trees built by the rule compiler never contain negated parameters.) (DEF 'NEGATION '[LAMBDA (ANT) (COND [(NULL ANT) NIL] [(EQUAL (CAR (EXPLODE ANT)) "#") (READLIST (CDR (EXPLODE ANT)))]

(READLIST (CDR (EXPLODE ANT)))] [T (READLIST (CONS '"#" (EXPLODE ANT)))])])

ECHO

Arguments	:	FILENAME the name of a file (in double quotes) which is to receive the text generated
		by INFER, IT or SET_TREES
Called by	:	the user.
Calls		builtin functions only
RETURNS	:	ECHO captures keyboard input and passes that
		input to EVAL. However, any text that is
		generated is also placed in FILENAME. This
		continues until the user requests that the

```
echo be turned off, or the user invokes a
             function which is not one of INFER, IT or
            SET TREES
(DEF 'ECHO
 [LAMBDA (FILENAME)
  [PROG (TEMP)
      (SETQ ECHO ON? T)
      (SETO IELOG
          (OUTPUT FILENAME))
      (PRINTC '"INFERENCE ENGINE LOG:" IELOG)
      (TERPRI NIL IELOG)
   LOOP
      (PRINTC '"+")
      (PRINTC '"+" IELOG)
      (SETQ TEMP
          (READ))
      (COND
          [(EQUAL TEMP
                'ECHO OFF)
             (SETQ ECHO ON? NIL)
             (PRINC '"END OF LOG." IELOG)
             (CLOSE IELOG)
             (TERPRI)
             (RETURN 'ECHO_OFF)]
          [(OR (ATOM TEMP)
                (EQUAL (CAR TEMP)
                    'INFER)
                (EQUAL (CAR TEMP)
                    'SET_TREES)
                (EQUAL (CAR TEMP)
                    'IT))
            (PRIN TEMP IELOG)
            (PRINT (PRIN (EVAL TEMP)) IELOG)
             (TERPRI)
            (TERPRI NIL IELOG)
            (GO LOOP)]
          ГΤ
            (SETQ ECHO ON? NIL)
            (PRINC '"END OF LOG." IELOG)
            (CLOSE IELOG)
            (PRIN (EVAL TEMP))
            (TERPRI)
            (TERPRI)
            (RETURN 'ECHO_OFF)])]])
                       DEL CURR GOAL
Arguments : LIST -- the list of parameter currently being
             inferred which is IM WORKING ON, ELE -- the
             parameter that has just been inferred
Called by :
             INFER2
Calls
          : builtin functions only
RETURNS : an updated IM WORKING ON list
```

```
(DEF 'DEL CURR_GOAL
'[LAMBDA (LIST ELE)
  (COND
     [(NULL LIST)
     NIL]
     [(EQUAL (CAR LIST) ELE)
        (CDR LIST)]
     [T
        (CONS (CAR LIST)
        (DEL_CURR_GOAL (CDR LIST) ELE))])])
```

К В

Arguments : NC -- number of complex rules, NS -- number of simple rules, NP -- number of parameters to be concluded about, FILENAME -- file in double quotes to send rules to when generated Called by : the user Calls : KNOWLEDGE BASE RETURNS : the knowledge base (DEF 'K B [LAMBDA (NC NS NP FILENAME HEURISTIC NUMB) [PROG (RESULTS) (SETQ ANT INFO (ARRAY 2 105 8))(SETQ LIST_OF_AVER_ANT NIL) (SETQ TL RULES 0) (SETQ LIST OF ASKFIRST '((a)(b)(c)(d)(e)(#a)(#b)(#c)(#d)(#e))) (SETO RESULTS (APPEND (KNOWLEDGE BASE NC NS NP 1) (KNOWLEDGE BASE 4 0 26 6))) (SETQ OP (OUTPUT FILENAME)) (PP RESULTS OP) (CLOSE OP) (RETURN RESULTS)]]) KNOWLEDGE BASE Arguments : N COMPLEX -- number of complex rules, N SIMPLE -- number of simple rules, N CONCL -number of conclusions, BEG -- place in alphabet to begin Called by : K_B : K B CREATOR Calls : the knowledge base RETURNS (DEF 'KNOWLEDGE BASE '[LAMBDA (N COMPLEX N SIMPLE N CONCL BEG) (K B CREATOR BEG N CONCL N COMPLEX N SIMPLE)]) K_B_CREATOR

Arguments : L -- parameter will be generating rules about, NUMBER OF PAR -- total number of conclusions, N COMPLEX -- number of complex rules, N SIMPLE -- number of simple rules Called by : KNOWLEDGE BASE

MEMBERS

Arguments : ELEMENT -- an antecedent to be tested, LST -the list being tested for ELEMENT CONSISTENCY Called by : CALLS MEMBER : RETURNS : T if ELEMENT is in LST, NIL otherwise (DEF 'MEMBERS '[LAMBDA (ELEMENT LST) (COND [(NULL (MEMBER ELEMENT LST)) NIL] [Т] т])])

GENERATE RULES

N ANT -- lower bound of interval of antecedent Arguments : L -- current conclusion, N RULES -- number of rules to generate for this conclusion Called by : K B CREATOR : RANDOM NUMBER, CREATE LIST, COUNT ANT Calls : the list of rules concluding about L RETURNS (DEF 'GENERATE RULES '[LAMBDA (N_ANT L N_RULES FLAG) $[PROG (I \overline{J} K M N)]$ (SETQ I 1) LOOP (SETQ J (FIX (+ N ANT (random number) 3)))) (SETQ K 1)(SETQ N NIL) (COND [(LE I N RULES) (SETQ M (CONS (CONS (SETQ LIST_OF_ANT

(CREATE LIST J K N L)) (LIST (READLIST (LIST (ASCII (+ L 64)))))) M)) (SETQ TL RULES (ADDI TL RULES)) (COUNT_ANT TE RULES $(+L 64) \overline{J} LIST OF ANT)$ (SETQ I (+ I 1))(GO LOOP)] ГТ (SETQ LIST OF AVER ANT (APPEND (LIST (CONS (READLIST (LIST (ASCII (+ L 64)))(LIST (FIND AVER ANT 1 N RULES)))) LIST OF AVER ANT)) (NOTE_LIST_OF_ANT_USED M) (SETQ LIST SO FAR NIL) (RETURN (REVERSE M))])]])

COUNT ANT

Arguments : ARRAY NUMBER -- array number in ANT INFO, CONCLUS -- current conclusion, MAX NO ANT -number of upper case letters in rule, LIST OF ANT -- list of antecedents in rule Called by : GENERATE RULES Calls NUMBER OF LOWER CASE : RETURNS NIL, the purpose is to store information in : ANT INFO (DEF 'COUNT ANT '[LAMBDA (ARRAY NUMBER CONCLUS MAX NO ANT LIST OF ANT) [PROG () (STORE (ANT INFO ARRAY NUMBER 1) CONCLUS) (STORE (ANT INFO ARRAY NUMBER '2) (NUMBER OF LOWER CASE LIST OF ANT 0)) (STORE (ANT INFO ARRAY NUMBER T3) MAX NO ANT) (COND [(GT (- MAX NO ANT (ANT INFO ARRAY NUMBER 2)) 0) (STORE REQ ANTS 1 (- MAX NO ANT (ANT INFO ARRAY NUMBER 2)) ARRAY NUMBER)]) (RETURN)]])

NUMBER_OF_LOWER_CASE

Arguments : LIST OF ANT -- list of antecedents in rule, NUMBER -- counter for the number of lower case

```
antecedents in rule
Called by : COUNT ANT
Calls
        : itself
RETURNS
        : the number of lower case antecedents in rule
(DEF 'NUMBER OF LOWER CASE
'[LAMBDA (LIST OF ANT NUMBER)
  (COND
      [(NULL LIST OF ANT)
        NUMBER]
      [(GE (CHRVAL (CAR LIST OF ANT)) 97)
        (NUMBER OF LOWER CASE (CDR LIST OF ANT)
            (ADD1 NUMBER)]
      [(EQUAL '"#"
           (CAR (EXPLODE (CAR LIST OF ANT))))
        (COND
            [(GE (CHRVAL (READLIST (CDR (EXPLODE (CAR
                                   LIST OF ANT)))) 97)
               (NUMBER OF LOWER CASE (CDR LIST OF ANT)
                   (AD\overline{D}1 \ \overline{N}UMBER)]
            [T]
              (NUMBER OF LOWER CASE (CDR LIST OF ANT)
                                                 NUMBER)])]
      ГΤ
        (NUMBER OF LOWER CASE (CDR LIST OF ANT) NUMBER)])])
                        CREATE LIST
Arguments : J -- number of antecedents for this rule,
             K -- counter of number of antecedents gener-
             ated so far, N -- NIL, L -- current conclu-
             sion
Called by : GENERATE RULES
            itself, CONSISTENCY, NUMBER REFORM
Calls
          :
          : the list of antecedents for a rule
RETURNS
(DEF 'CREATE LIST
'[LAMBDA (J K N L)
  (COND
      [(LE K J)
        (CONS (SETQ TEMP
                   (CONSISTENCY L N))
            (CREATE LIST J
                 (+1 K)
                 (CONS (CHRVAL (NUMBER REFORM TEMP)) N) L))]
      [т]
        NIL])])
```

NUMBER REFORM

Arguments : TEMP -- an antecedent Called by : CREATE_LIST

```
: LISP buildin functions
Calls
RETURNS
             its purpose is to strip the "not" symbol from
          :
             an antecedent to be used in the list of ants
             generated for the rule so far
(DEF 'NUMBER REFORM
'[LAMBDA (TEMP)
  (COND
      [(EQUAL '"#"
           (CAR (EXPLODE TEMP)))
        (READLIST (CDR (EXPLODE TEMP)))]
      Гт
        TEMP])])
                        CONSISTENCY
Arguments : L -- the current conclusion, N -- the list of
             antecedents generated so far
Called by :
             CREATE LIST
             MEMBERS, TEST FOR NOT, LOOK AT ANTS USED BY
Calls
          :
RETURNS
          : an antecedent
(DEF 'CONSISTENCY
'[LAMBDA (L N)
  [PROG (CONCL)
      (SETQ CONCL
          (+ 64 L))
   LOOP
      (SETQ PRELIM
          (FIX (+ 56
                    (* (RANDOM NUMBER) 35))))
      (COND
          [(EQUAL PRELIM CONCL)
            (GO LOOP)]
          [(MEMBERS PRELIM N)
            (GO LOOP)]
          [(LT PRELIM 70)
            (SETQ PRELIM
                (+ PRELIM 41))
            (COND
                [(MEMBERS PRELIM N)
                  (GO LOOP)]
                ГΤ
                       (RETURN (TEST FOR NOT PRELIM))])]
          [(LE PRELIM CONCL)
            (COND
                [(EQUAL (LOOK AT ANTS USED BY (READLIST
                    (LIST (ASCII PRELIM)))
                           (READLIST (LIST (ASCII CONCL))))
                      'т)
                   (GO LOOP)]
                Гт
                   (RETURN (TEST FOR NOT PRELIM))])]
```

[T (RETURN (TEST_FOR_NOT PRELIM))])]]) TEST_FOR_NOT Arguments : PRELIM -- the antecedent just generated Called by : CONSISTENCY Called by : LICE built in functions

Calls : LISP builtin functions RETURNS : a negated antecedent averaging 1 in 10 DEF 'TEST_FOR_NOT [LAMBDA (PRELIM) (COND [(EQUAL (FIX (+ 1 (* (RANDOM_NUMBER) 10))) 1) (SETQ PRELIM (READLIST (CONS '"#" (EXPLODE (READLIST (LIST (ASCII PRELIM)))))))] [T (SETQ PRELIM

(READLIST (LIST (ASCII PRELIM))))])])

PRINT STATS

Arguments : FILENAME -- filename, in double quotes, of file that will contain the information in ANT INFO Called by : the user Calls : PRIN ARR RETURNS : outputs the contents of ANT_INFO (DEF 'PRINT STATS '[LAMBDA (FILENAME) [PROG (K) (SETO OF (OUTPUT FILENAME)) (SETQ K 1)LOOP (COND [(LT K 105) (PRINT ARR 1 7 K) $(SETQ \overline{K})$ (ADD1 K)) (GO LOOP)] ГΤ (CLOSE OF)]) (RETURN)]])

PRIN ARR

INDEX -- counter, MAX -- number of columns, Arguments : K -- current row in array Called by : PRINT STATS Calls itself, builtin functions : RETURNS prints elements of array ANT INFO to a file : (DEF 'PRIN ARR '[LAMBDA (INDEX MAX K) (COND [(LE INDEX MAX) (SETQ ANS (ANT INFO K INDEX)) $(PRIN AN\overline{S} OF)$ (SPACES 3 OF) (PRINT ARR (ADD1 INDEX) MAX K)] ſТ (TERPRI NIL OF)])]) LOOK AT ANTS USED BY Arguments : ANT CANDIDATE -- antecedent candidate, CURR CONCL -- conclusion of the rule the antecedent is being tested for Called by : CONSISTENCY FIND CORRECT ANT LIST Calls : : T if the current conclusion appears as an RETURNS antecedent in a rule concluding about the ANT CANDIDATE, NIL otherwise (DEF 'LOOK AT ANTS USED BY '[LAMBDA (ANT_CANDIDATE CUR_CONCL) (FIND CORRECT ANT LIST ANT CANDIDATE CUR_CONCL LIST OF ANT & THEIR CONCL)]) FIND CORRECT_ANT_LIST Arguments : ANT CANDIDATE -- antecedent candidate, CUR CONCL -- conclusion of current rule, LIST --list of conclusions and the antecedents used by them LOOK_AT_ANTS_USED_BY Called by : itself, builtin functions Calls : T if CUR CONCL appears as an antecedent in a RETURNS : rule concluding about ANT_CANDIDATE, NIL otherwise (DEF 'FIND CORRECT ANT_LIST '[LAMBDA (ANT CANDIDATE CUR_CONCL LIST) (COND [(NULL LIST)

NIL] [(EQUAL (CAAR LIST) ANT_CANDIDATE) (COND [(NULL (MEMBER CUR_CONCL (CDAR LIST))) NIL] [T [T (FIND_CORRECT_ANT_LIST_ANT_CANDIDATE_CUR_CONCL (CDR LIST))])]

NOTE LIST OF ANT_USED

MAKE LIST FROM RULES

Calls	:	RULELIST NOTE_LIST_OF_ANT_USED PASS_ANT_LIST, itself the list of antecedents us	sed	in	the	rules
(DEF 'MAKE '[LAMBDA ((COND		IST_FROM_RULES LELIST)				
[(NU	JLL	RULELIST)				
LI	ST	SO FAR]				
[Т	-					
		S_ANT_LIST (CAAR RULELIST)) C_LIST_FROM_RULES (CDR RULE		ST)))

PASS_ANT_LIST

Arguments	:	ANTLIST list of antecedents of one rule
Called by	:	MAKE_LIST FROM RULES
Calls	:	ANT TO CHECK, itself
RETURNS	:	an updated list of antecedents

ANT TO CHECK

```
ANT -- antecedent to check
Arguments :
Called by : PASS ANT LIST
Calls
             MEMBER
          :
RETURNS
          :
             a list of antecedents with ANT added to it if
             it is not already present. The list contains
             antecedents with no negation symbols
(DEF 'ANT TO CHECK
'[LAMBDA (ANT)
  (COND
      [(EQUAL '"#"
           (CAR (EXPLODE ANT)))
        (COND
            [(NULL (MEMBER (READLIST (CDR (EXPLODE ANT)))
                   LIST_SO_FAR))
               (SETQ ANT TO ADD
                   (READLIST (CDR (EXPLODE ANT))))
              T]
            [Т]
              NIL])]
      ГΤ
        (COND
            [(NULL (MEMBER ANT LIST_SO_FAR))
              (SETQ ANT TO ADD ANT)
              T]
            [T]
              NIL])])])
                      RANDOM NUMBER
Arguments :
            none
Called by : CONSISTENCY
          : builtin functions
Calls
RETURNS
            a random number between 0 and 1
          :
(DEF 'RANDOM NUMBER
'[LAMBDA ()
 [PROG ()
```

FIND ALL

Arguments : ANT -- an antecedent, RULESET -- a rule partition Called by : COMPILE, itself ANT IS IN RULE, itself Calls : : the list of all rules which have ANT as an RETURNS antecedent. (DEF 'FIND ALL '[LAMBDA (ANT RULESET) (COND [(NULL RULESET) NIL] [(ANT_IS_IN_RULE ANT (CAR RULESET)) (CONS (CAR RULESET) (FIND ALL ANT (CDR RULESET)))] ГΤ (FIND ALL ANT (CDR RULESET))])]) ANT IS IN RULE ANT -- an antecedent, RULE -- a rule Arguments : Called by : FIND ALL ANT IS IN LIST Calls : Predicate - T if the antecedent is present in RETURNS : in the rule, NIL otherwise. (DEF 'ANT IS IN RULE '[LAMBDA (ANT RULE) (ANT IS IN LIST ANT $\overline{(CAR RULE)}$ ANT IS IN LIST Arguments : ANT -- an antecedent, LIST_ANT -- a list of antecedents Called by : ANT_IS_IN_RULE, itself : itself Calls T if the antecedent is in the list RETURNS : Predicate. of antecedents, NIL otherwise.

(DEF 'ANT_IS_IN_LIST '[LAMBDA (ANT LIST_ANT) (COND [(NULL LIST_ANT) [(EQUAL (CAR LIST_ANT) ANT) T] [T (ANT_IS_IN_LIST_ANT (CDR_LIST_ANT))])])

NIL]

DEL RULE

Arguments : ANT -- an antecedent, RULE -- a rule DEL RULESET Called by : Calls DEL ANT LIST : RETURNS The rule which is RULE without antecedent ANT : (DEF 'DEL RULE '[LAMBDA (ANT RULE) (COND [(NULL RULE) NIL] [Т] (CONS (DEL ANT LIST ANT (CAR RULE)) (CDR RULE))])]) DEL ANT LIST Arguments : ANT -- an antecedent, ALIST -- an antecedent list DEL RULE, itself Called by : Calls itself : the antecedent list which is ALIST with all of RETURNS : the occurrences of ANT removed. (DEF 'DEL ANT_LIST '[LAMBDA (ANT ALIST) (COND [(NULL ALIST) NIL] [(EQUAL (CAR ALIST) ANT) (DEL ANT LIST ANT (CDR ALIST))] ГΤ (CONS (CAR ALIST) (DEL ANT LIST ANT (CDR ALIST)))])])

DEL RULESET Arguments : ANT -- an antecedent, RULESET -- a rule partition (list of rules) Called by : COMPILE, itself Calls : DEL_RULE, itself RETURNS : the rule set which is RULESET with ANT removed from the antecedent list of each rule. (DEF 'DEL RULESET '[LAMBDA (ANT RULESET) (COND [(NULL RULESET) NIL] ĽТ (CONS (DEL RULE ANT (CAR RULESET)) (DEL RULESET ANT (CDR RULESET)))])])

PACKAGE "HEURS1.LSP"

SELECT_A_CLAUSE

Arguments	:	RULESET a rule partition, or list of rules concluding about the same parameter
Called by	:	-
Calls	:	MOST OFTEN OCCURRING, RULE_GROUP_TALLY
	•	an antecedent which will be placed nearest
	•	the top of a decision tree. In this package the heuristics used is the antecedent that appears in the most rules will be placed nearest the root of the tree.
(DEF 'SELE '[LAMBDA (
[PROG ()	ŀ	
(SET	Q :	SET
(CHEC	K I	RULESET_FOR_ASKFIRST LIST_OF_ASKFIRST RULESET))
(CON	D_	
·	F (1	NULL SET)
		(RETURN (MOST_OFTEN_OCCURRING (RULE_GROUP_TALLY RULESET)))]
	[Т]	
(RE	TUI	RN (MOST_OFTEN_OCCURRING (RULE_GROUP_TALLY SET)

MOST_OFTEN_OCCURRING

Called by	:	TALLYLIST a list of antecedent/count pairs SELECT_A_CLAUSE
Calls	:	MOST
RETURNS	:	the antecedent whose count in tallylist is the greatest. If two antecedents tie, the "leftmost" one in TALLYLIST is returned

(DEF 'MOST_OFTEN_OCCURRING '[LAMBDA (TALLYLIST) (MOST_TALLYLIST_NIL_0)])

MOST

Arguments	•	TLIST a tally list, CLAUSE the clause which has been determined to be the most often occurring (so far), COUNT the count of this
Calle	•	parameter MOST_OFTEN_OCCURRING, itself itself the antecedent with the highest count. MOST does the work described under MOST_OFTEN_OCCUR.

TALLY

Arguments : ANTECEDENT -- an antecedent which we want to update tally for, TALLYLIST -- a list of antecedent/count pairs Called by : LIST TALLY, itself Calls : NEGATION, itself RETURNS : a tally list which contains an updated entry for ANTECEDENT. If there was no entry for ANTECEDENT, then one is built with a count of l. If a negated antecedent is found, the entry is updated for the unnegated antecedent

(DEF 'TALLY

'[LAMBDA (ANTECEDENT TALLYLIST) (COND [(NULL TALLYLIST) (LIST (LIST ANTECEDENT 1))] [(EQUAL (CAAR TALLYLIST) ANTECEDENT) (CONS (LIST ANTECEDENT (ADD1 (CADAR TALLYLIST))) (CDR TALLYLIST))] [(EQUAL (CAAR TALLYLIST) (NEGATION ANTECEDENT)) (CONS (LIST ANTECEDENT (ADD1 (CADAR TALLYLIST))))] Γт (CONS (CAR TALLYLIST) (TALLY ANTECEDENT (CDR TALLYLIST)))])])

LIST TALLY

Arguments	:	ANT LIST a list of antecedents, TLIST a
		tally list
Called by	:	RULE_TALLY, itself
Calls		TALLY, itself
RETURNS	:	performs the TALLY function for each
		antecedent in ANT_LIST

(DEF 'LIST_TALLY '[LAMBDA (ANT_LIST TLIST) (COND [(NULL ANT_LIST) TLIST] [T (LIST_TALLY (CDR ANT_LIST) (TALLY (CAR ANT_LIST) TLIST))])])

RULE TALLY

Arguments : RULE -- a rule of the form (antecedent-list consequent), TLIST -- a tally list Called by : RULE_GROUP_TALLY Calls : LIST_TALLY : calls LIST TALLY to process its antecedent RETURNS list (DEF 'RULE TALLY '[LAMBDA (RULE TLIST) (COND [(NULL RULE) RULE] [Т] (LIST TALLY (CAR RULE) TLIST)])]) RULE GROUP TALLY Arguments : R GROUP -- a rule group (a list of rules) Called by : SELECT A CLAUSE, itself RULE)TALLY, itself Calls : a tally list of all the antecedents used in RETURNS : R GROUP. This function processes each rule by successively calling RULE TALLY for each rule in R GROUP (DEF 'RULE GROUP TALLY $[LAMBDA (\overline{R} GROUP)]$ (COND [(NULL R GROUP) NIL] ſт (RULE TALLY (CAR R GROUP) (RULE_GROUP_TALLY (CDR R_GROUP)))])]) CHECK ANT LIST Arguments : AF -- a list of known parameters (askfirst), ANT LIST -- a list of antecedents of a rule

Called by :

:

Calls

ANT IS IN LIST

LOOK_AT_RULE

'[LAMBDA (ĀF RULE) (CHECK ANT LIST AF (CAR RULE))])

CHECK RULESET FOR ASKFIRST

AF -- list of known parameters, RULESET -- a Arguments : list of rules Called by : : LOOK_AT_RULE, itself Calls RETURNS : (DEF 'CHECK RULESET FOR ASKFIRST '[LAMBDA (AF RULESET)](COND [(NULL RULESET) NIL] [(LOOK AT RULE AF (CAR RULESET)) (CONS (CAR RULESET) (CHECK_RULESET_FOR_ASKFIRST AF $(C\overline{D}R RULES\overline{E}T))$ [Т] (CHECK RULESET FOR ASKFIRST AF (CDR RULESET))])

SELECT A CLAUSE

Arguments : RULESET -- a partition list Called by : COMPILE Calls : TRAVEL LIST MIN ANT : an antecedent to be put out as a branch RETURNS (DEF 'SELECT A CLAUSE '[LAMBDA (RULESET) (COND [(NULL (SETQ CLAUSE (TRAVEL_LIST MIN ANT RULESET LIST OF ASKFIRST))) (TRAVEL_LIST MIN ANT RULESET LIST OF MIN ANT)] Гт CLAUSE])]) CREATE AVERAGES Arguments : none Called by : K B REORDER, NEW AVERAGES, ADD NOTS Calls : : an updated LIST OF AVER ANT after calculating RETURNS averages three times (DEF 'CREATE AVERAGES '[LAMBDA () [PROG () (SETQ LIST OF ANTS NIL) (SETQ AVER ONE (REORDER 0 4 (LIST OF AVER ANT))) (SETQ LIST OF ANTS NIL) (SETQ AVER TWO (REORDER 0 16

(NEW_AVERAGES 6 26 AVER_ONE))) (SETQ LIST_OF_ANTS NIL) (SETQ AVER_THREE

(REORDER 0 64 (NEW_AVERAGES 6 26 AVER_TWO))) (RETURN (ADD_NOTS AVER_THREE))]])

NEW AVERAGES

Arguments : ST -- letter to begin calculating averages, FIN -- letter to stop calculating averages (Z) AVER_LIST -- current list of average antecedents Called by : CREATE_AVERAGES Calls : EVAL_FOR_ALL_RULES, itself RETURNS : an updated list of averages (DEF 'NEW AVERAGES '[LAMBDA (ST FIN AVER_LIST) (COND [(GT ST FIN) NIL] Ет (SETQ NEW SUM & AVER (EVAL FOR ALL RULES 1 N RULES ST O AVER LIST)) (COND [(GE (% NEW_SUM_&_AVER N_RULES) 2) (SETQ NEW SUM & AVER (+ (FIX (/ NEW SUM & AVER N RULES)) 1))]ГΤ (SETQ NEW SUM & AVER (FIX (/ NEW SUM & AVER N_RULES)))]) (SETQ NEW AVER (CONS (CONS (READLIST (LIST (ASCII (+ ST 64)))) (LIST NEW SUM & AVER)) (NEW AVERAGES (ADD ST FIN AVER LIST)))])])

EVAL FOR ALL RULES

Arguments : IND -- index, NUMB RULES -- number of rules to evaluate, BLK -- previous block of rules in ANT INFO, R SUM -- sum of averages for rules, AVER LIST -- current list of averages NEW AVERAGES Called by : CALC NEW NUMB, itself Calls : : the new sum of antecedents RETURNS (DEF 'EVAL FOR ALL RULES '[LAMBDA (IND NUMB RULES BLK R_SUM AVER_LIST) (COND [(GT IND NUMB RULES) R_SUM] [Т] (COND [(EQUAL (SETQ N ANTS (- (ANT_INFO (+ IND BLK) 3) (ANT_INFO (+ IND BLK) 2))) 0) (EVAL_FOR_ALL_RULES (ADD1 IND) NUMB_RULES BLK R SUM AVER LIST)] [Т] (SETQ R SUM $(+ \overline{R} SUM)$ (CALC_NEW_NUMB (+ IND BLK) N_ANTS 0 AVER LIST)) (EVAL_FOR_ALL_RULES (ADD1 IND) NUMB_RULES BLK R_SUM AVER_LIST)])])

CALC NEW NUMB

Arguments : ARR NO -- array number, N ANTS -- number of antecedents that have to be inferred, ANT SUM -- sum of antecedents, AVER LIST -- current average of antecedents EVAL FOR ALL RULES Called by : GET NUMB OF NEEDED ANTS, itself : Calls RETURNS : new sum of antecedents (DEF 'CALC NEW NUMB '[LAMBDA (ARR NO N ANTS ANT_SUM AVER_LIST) (COND [(GT N ANTS O) (SETQ ANT SUM (+ (GET_NUMB_OF_NEEDED_ANT (ANT_INFO ARR_NO (+ 3 N ANTS))AVER LIST) ANT SUM)) (CALC NEW NUMB ARR NO (SUB1 N_ANTS) ANT_SUM AVER_LIST)] [т ANT SUM])])

REORDER

Arguments :: ST -- lowest average in list of antecedents, FIN -- highest average in list of antecedents, LIST -- current list of averages Called by : CREATE_AVERAGES PICK OUT : Calls a list of averages in ascending order RETURNS : (DEF 'REORDER '[LAMBDA (ST FIN LIST) (COND [(LE ST FIN) (PICK_OUT ST LIST) (REORDER (ADD1 ST) FIN LIST)] Гт LIST OF ANTS])])

PICK OUT

Arguments Called by Calls RETURNS	of averages
(DEF 'PICK '[LAMBDA ((COND	OUT THESE_CONCL LIST)

[(NULL LIST) LIST_OF_ANTS] [(EQUAL (CADAR LIST) THESE_CONCL) (SETQ_LIST_OF_ANTS (APPEND_LIST_OF_ANTS (LIST (CAR_LIST)))) (PICK_OUT_THESE_CONCL (CDR_LIST))] [T (PICK_OUT_THESE_CONCL (CDR_LIST))])])

STORE REQ ANTS

ST -- counter, N ANTS -- number of antecedents Arguments : to be stored, ARR NUMB -- array number Called by : Calls : WHAT THE ANT IS, itself nothing, the purpose is to store information RETURNS : about the rules in ANT_INFO (DEF 'STORE REQ ANTS '[LAMBDA (ST N ANTS ARR NUMB) (COND [(LE ST N ANTS) (STORE (ANT_INFO ARR_NUMB (+ st 3)) (WHAT_THE ANT_IS LIST_OF_ANT ST 1)) (STORE_REQ_ANTS (ADD1 ST) N_ANTS ARR_NUMB)])])

WHAT THE ANT IS

Arguments : LIST -- list of antecedents, N_TO_FIND-- number of antecedents to store, COUNT -- index Called by : STORE REQ_ANTS Calls : itself RETURNS : an antecedent (DEF 'WHAT THE ANT IS '[LAMBDA (LIST N_TO_FIND COUNT) (COND [(NULL LIST) NIL] [(EQUAL '"#" (CAR (EXPLODE (CAR LIST)))) (COND [(LT (CHRVAL (READLIST (CDR (EXPLODE (CAR LIST)))) 97) (COND [(EQUAL COUNT N TO FIND) (CHRVAL (READLIST (CDR (EXPLODE (CAR LIST)))))]

Гт (WHAT_THE_ANT IS (CDR LIST) N TO FIND (ADD1 COUNT))])] ГΤ (WHAT_THE_ANT_IS (CDR LIST N_TO_FIND COUNT))])] [Т] (COND [(LT (CHRVAL (CAR LIST)) 97) (COND [(EQUAL COUNT N TO FIND) (CHRVAL (CAR LIST))] Γт (WHAT THE ANT IS (CDR LIST) N TO FIND (ADD1 COUNT))])] [Т] (WHAT THE ANT IS (CDR LIST) N TO FIND COUNT)])]) FIND AVER ANT INDEX -- index to array number, N RULES--Arguments : number of rules for a particular conclusion Called by : Calls builtin functions : RETURNS new average : (DEF 'FIND AVER ANT '[LAMBDA (INDEX N RULES) [PROG () (SETQ SUM 0) LOOP (COND [(LE INDEX N RULES) (SETQ SUM (+ SUM (- (ANT_INFO (+ (- TL_RULES INDEX) 1) 3) (ANT INFO (+ (- TL_RULES INDEX) 1) 2))) (SETO INDEX (ADD1 INDEX)) (GO LOOP)] [T] (COND [(GE (% SUM N RULES) 2) (RETURN (+ (FIX (/ SUM N_RULES)) 1))] ſт (RETURN (FIX (/ SUM N_RULES)))])]))

ADD NOTS

Arguments : ORDERED_LIST -- list of conclusions and averages Called by : CREATE_AVERAGES

Calls itself : RETURNS the ordered list with negations : (DEF 'ADD NOTS '[LAMBDA (ORDERED LIST) (COND [(NULL (CAR ORDERED LIST)) NIL] [T] (APPEND (APPEND (LIST (CAR ORDERED LIST)) (LIST (CONS (READLIST (CONS '"#" (EXPLODE (CAAR ORDERED LIST)))) (LIST (CADAR ORDERED LIST))))) (ADD NOTS (CDR ORDERED LIST)))])]) GET NUMB OF NEEDED ANT Arguments : ANT -- current antecedent, LIST -- list of antecedents and their averages Called by : CALC NEW NUMB Calls itself : RETURNS : the average for ANT (DEF 'GET NUMB OF NEEDED ANT '[LAMBDA (ANT LIST) (COND [(NULL LIST) NIL] [(GE ANT 97) 0] [(EQUAL ANT (CHRVAL (CAAR LIST))) (CADAR LIST)] ГΤ (GET NUMB OF NEEDED ANT ANT (CDR LIST)))))) TRAVEL LIST_MIN_ANT Arguments : R_GROUP -- a partition list, LIST_ANTS -- list of antecedents and their averages SELECT A CLAUSE Called by : LOOK_FOR_ANT_IN_RULESET, itself Calls : RETURNS nothing : (DEF 'TRAVEL LIST MIN ANT '[LAMBDA (R GROUP LIST ANTS) (COND [(NULL LIST_ANTS) [(LOOK_FOR_ANT_IN_RULESET (CAAR LIST_ANTS) R_GROUP) (CAAR LIST ANTS)]

[T (TRAVEL LIST MIN ANT R GROUP (CDR LIST_ANTS))])])

LOOK_FOR_ANT_IN_RULESET

Arguments : ANT -- an antecedent, R_GROUP -- a partition list TRAVEL LIST MIN ANT Called by : ONE RULE, itself Calls : T if antecedent is found, NIL otherwise RETURNS : (DEF 'LOOK FOR ANT IN RULESET '[LAMBDA (ANT R GROUP) (COND [(NULL R_GROUP) NIL] [(ONE RULE ANT (CAR R GROUP)) T] [Т] (LOOK FOR ANT IN RULESET ANT $(\overline{C}DR \overline{R} GR\overline{O}UP)))))$ ONE_RULE ANT -- antecedent, RULE -- a rule Arguments : LOOK FOR ANT IN RULESET Called by : itself Calls : T if a match is found between an antecedent RETURNS : in a rule and in the list of antecedents (DEF 'ONE RULE '[LAMBDA (ANT RULE) (COND [(MEMBER ANT (CAR RULE)) т] [T] NIL])])